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## On Automatic Synthesis of Analog/Digital Circuits

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**Abstract**— The paper builds on a recent explicit numerical algorithm for Kolmogorov's superpositions, and will show that in order to synthesise minimum *size* (*i.e.*, *size*optimal) circuits for implementing any Boolean function, the nonlinear activation function of the gates has to be the identity function. Because classical AND-OR implementations, as well as threshold gate implementations require exponential *size*, it follows that *size*-optimal solutions for implementing arbitrary Boolean functions can be obtained using analog (or mixed analog/digital) circuits. Conclusions and several comments are ending the paper.

#### I. INTRODUCTION

In this paper a *network* is an acyclic graph having several input nodes, and some output nodes. If a synaptic *weight* is associated with each edge, and each node computes the weighted sum of its inputs to which a nonlinear activation function is applied:  $f(x_1, ..., x_{\Delta}) = \sigma(\sum_{i=1}^{\Delta} w_i x_i + \theta)$ , the network is a *neural network* (NN), with  $w_i \in \mathbb{R}$  the synaptic weights,  $\theta \in \mathbb{R}$  known as the *threshold*,  $\Delta$  being the *fan-in*, and  $\sigma$  a non-linear activation function. The underlying graph is acyclic, thus the network can be layered, *i.e.*, it is a *multilayer feedforward NN*. It is characterised by: *depth* (*i.e.*, number of layers), and *size* (*i.e.*, number of neurons).

The paper starts by overviewing results dealing with the approximation capabilities of NNs, and known upper and lower bounds on the *size* of *threshold gate circuits* (TGCs). They show that TGCs require exponential *size* for implementing arbitrary *Boolean functions* (BFs). Based on a constructive solution for Kolmogorov's superpositions we will prove that in order to obtain linear *size* NNs (*i.e., size-optimal*) for implementing any BF, *the nonlinear activation function of the neurons has to be the identity function*. Because both Boolean and TG circuits require exponential *size*, it follows that *size-optimal* implementations of BFs can be obtained in analog circuitry. Conclusions, and comments on the required precision are ending the paper.

#### **II. PREVIOUS RESULTS**

NNs have been experimentally shown to be quite effective in many applications (see Applications of Neural Networks in [2], together with Part F: Applications of Neural Computation and Part G: Neural Networks in Practice: Case Studies from [20]). This success generated two directions of research for finding: (i) existence/constructive proofs for the 'universal approximation problem'; (ii) tight bounds on the size needed by the approximation problem. The paper will focus on both aspects, for the case when the functions to be implemented are BFs.

#### A. Neural Networks as Universal Approximators

The first line of research has concentrated on the approximation capabilities of NNs [14,22,35,36]. It was started in 1987 by Hecht-Nielsen [26] and Lippmann [47] who, together with LeCun [45], were probably the first to recognise that the specific format of Kolmogorov's superpositions [41]  $f(x_1,...,x_n) = \sum_{q=1}^{2n+1} \Phi_q(y_q)$  in [67,68]:

$$f(x_1,...,x_n) = \sum_{q=1}^{2n+1} \left\{ \Phi_q \left[ \sum_{p=1}^n \alpha_p \, \psi(x_p + qa) \right] \right\}$$
(1)

can be interpreted as a NN with one hidden layer. This gave an existence proof of the approximation properties of NNs. The first nonconstructive proof was given in 1988 by Cybenko [16,17], and was independently presented by Irie&Miyake [34]. Similar results for radial basis functions were shortly reported [24,59]. Different enhancements have been later presented (see [9,64]):

- Funahashi [21] proved the same result and also refined the use of Kolmogorov's theorem in [26], giving an approximation result for two-hidden-layer NNs;
- Hornik et al. [31] showed that the continuity requirement for the output function can partly be removed;
- Hornik et al. [32] also proved that a NN can approximate simultaneously a function and its derivative;
- Park&Sandberg [57, 58] used radial basis functions in the hidden layer (almost constructive proof);
- Hornik [29] showed that the continuity requirement can be completely removed, the activation function having to be 'bounded and nonconstant';
- Geva&Sitte [23] proved that four-layered NNs with sigmoid activation are universal approximators;
- Kůrková [43,44] has demonstrated the existence of approximate superposition representations within the constraints of NNs, *i.e.*  $\psi$  and  $\Phi_q$  being of the form  $\sum a_r \sigma (b_r x + c_r)$ , where  $\sigma$  is an arbitrary activation sigmoidal function;
- Mhaskar&Micchelli [49,50] approach was based on the Fourier series of the function, by truncating the infinite sum, and rewriting  $e^{ikx}$  in terms of the activation function (which now has to be periodic);
- Koiran [40] presented a proof on the line of Funahashi's proof [21], which allows the use of units with 'piecewise continuous' activation functions;
- Leshno *et al.* [46] relaxed the condition for the activation function to 'locally bounded piecewise continuous', thus embedding as special cases almost all the activation functions that have been reported;
- Hornik [30] added to these results by proving that: (i) if the activation function is locally Riemann integrable and nonpolynomial, the *weights* and the *thresholds* can be constrained to arbitrarily small sets; and (ii) if the activation function is locally analytic, a single universal *threshold* will do;
- Funahashi&Nakamura [22] showed that the universal approximation theorem also holds for trajectories;

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- Sprecher [69] has demonstrated that there are universal hidden layers that are independent of n;
- Barron [4] described spaces of functions that can be approximated by the relaxed algorithm of Jones [37] using functions computed by single-hidden-layer networks of perceptrons;
- Attali&Pagès [3] provide an elementary proof based on the Taylor expansion and the Vandermonde determinant, yielding bounds for the design of the hidden layer and convergence results for the derivatives.

All these results were obtained provided that sufficiently many hidden units are available. More constructive solutions have been obtained in very small depth [38,54,55]. but their size still grows fast with respect to the number of dimensions and/or examples, or with the required precision. Recently, an explicit numerical algorithm for superpositions has been detailed [70-72].

#### **B. Threshold Gate Circuits**

The other line of research was to find the smallest size NN which can realise an arbitrary function given a set of *m* vectors from  $\mathbb{IR}^n$ . Many results have been obtained for TGs [51]. The first lower bound of:

$$size \ge 2(2^{n}/n)^{1/2}$$
 (2)

on the size of a TGC for 'almost all' n-ary BFs (i.e.,  $f: \mathbb{B}^n \to \mathbb{B}$ ) was given in [53]. Later [48] a very tight upper bound was proven in depth = 4:

size 
$$\leq 2 (2^{n}/n)^{1/2} \times \{1 + \Omega [(2^{n}/n)^{1/2}]\}.$$
 (3)

A similar existence exponential lower bound of  $\Omega(2^{n/3})$ for arbitrary BFs can be found in [65] (see also [63]).

For classification problems, the first result was that a NN of depth = 3 and size = m - 1 could compute an arbitrary dichotomy. The main improvements have been:

- Baum [5] presented a TGC with one hidden layer having  $\lceil m/n \rceil$  neurons realising an arbitrary dichotomy on a set of m points in general position in  $\mathbb{IR}^{n}$ ; if the points are on the corners of the n-dimensional hypercube ( $f: \mathbb{IB}^n \to \mathbb{IB}$ ), m-1 nodes are still needed;
- a tighter bound of [1 + (m-2)/n] neurons in the hidden layer for realising an arbitrary dichotomy on a set of *m* points which satisfy a more relaxed topological assumption was proven in [33]; the m-1 nodes condition was shown to be the least upper bound needed;
- Arai [1] showed that m-1 hidden neurons are necessary for arbitrary separability, but improved the bound for the dichotomy problem to m/3;
- Beiu&De Pauw [10] have detailed existence lower and upper bounds:  $2m / (n \log n) < size < 1.44m / n$  by estimating the entropy of the data-set ([11,13]).

Other existence lower bounds for the arbitrary dichotomy problem can be found in [25,59]:

- a depth-2 TGC requires m/{n log(m/n)} TGs;
  a depth-3 TGC requires 2 (m/logm)<sup>1/2</sup> TGs in each of the two hidden layer (if  $m \gg n^2$ );
- an arbitrarily interconnected TGC without feedback needs  $(2m/\log m)^{1/2}$  TGs (if  $m \gg n^2$ ).

One study [15] has tried to unify these two lines of research by first presenting analytical solutions for the general NN problem in one dimension (having infinite size), and then giving practical solutions for the one-dimensional cases (i.e., including an upper bound on the size). Extensions to the *n*-dimensional case using three- and four-layers solutions were derived under piecewise constant approximations, and under piecewise linear approximations.

#### **C. Boolean Functions**

The particular case of BFs has been studied intensively [9,56]. Many results have been obtained for particular BFs [63,65], but a size-optimal result for BFs that have mgroups of ones in their truth table  $IF_{n,m}$  (i.e., BFs defined on the *m* groups) was detailed by Red'kin in 1970 [62].

Theorem from [62] The complexity realisation (i.e., number of elements) of  $IF_{n,m}$  is at most 2  $(2m)^{1/2} + 3$ .

This result is valid for unlimited fan-in TGs. Departing from these lines, Horne&Hush [28] have detailed a solution for limited fan-in TGCs.

**Theorem from [28]** Arbitrary BFs  $f: \{0, 1\}^n \rightarrow \{0, 1\}^m$ can be implemented in a NN of perceptrons restricted to fan-in 2 with a node complexity of  $\Theta$  {m 2<sup>n</sup>/(n + logm)} and requiring O(n) layers.

### **II. SIZE-OPTIMAL IMPLEMENTATIONS**

Implementing arbitrary BFs using classical Boolean gates (i.e., AND and OR) requires exponential size. The known bounds for size are also exponential if TGCs are used for solving arbitrary BFs [6]. These bounds reveal exponential gaps, and also suggest that TGCs with more layers (depth  $\neq$  constant [8,12]) might have a smaller size.

Another approach is to use Kolmogorov's superpositions, which shows that there are NNs having only 2n + 1neurons which can approximate any function. Such a solution is clearly size-optimal. We start from [70-72], where a constructive solution for the general case was detailed.

**Theorem from [70]** Define the function  $\Psi : \mathcal{E} \to \mathcal{D}$  such that for each integer  $k \in N$ :

$$\Psi\left(\sum_{r=1}^{k} i_{r} \gamma^{-r}\right) = \sum_{r=1}^{k} \tilde{i}_{r} 2^{-m_{r}} \gamma^{-\frac{n^{r-m_{r-1}}}{n-1}}$$
(4)

where  $i_r = i_r - (\gamma - 2) \langle i_r \rangle$  and

$$m_{r} = \langle i_{r} \rangle \left\{ 1 + \sum_{s=1}^{r-1} [i_{s}] \times \dots \times [i_{r-1}] \right\}$$
(5)

for r = 1, 2, ..., k.

Here  $\gamma \ge 2n+2$  is a base,  $\mathcal{E} = [0, 1]$  is the unit interval.  $\mathcal{D} \text{ is the set of rational numbers } d_k = \sum_{r=1}^k i_r \gamma^{-r} \text{ defined on } k \in N \text{ digits } (0 \le i_r \le \gamma - 1). \text{ Also, } \langle i_1 \rangle = [i_1] = 0, \text{ while for } r \ge 2; \langle i_r \rangle = 0 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2, \langle i_r \rangle = 1 \text{ when } i_r = 0, 1, \dots, \gamma - 2 \text{ when } i_r = 0, 1, \dots, \gamma - 2 \text{ when } i_r = 0, 1, \dots, \gamma - 2 \text{ when } i_r = 0, 1, \dots, \gamma - 2 \text{ when } i_r = 0, 1, \dots, \gamma - 2 \text{ when } i_r = 0, 1, \dots, \gamma - 2$  $i_r = \gamma - 1$ ,  $[i_r] = 0$  when  $i_r = 0, 1, ..., \gamma - 3$ , while  $[i_r] = 1$ when  $i_r = \gamma - 2$ ,  $\gamma - 1$ . If we limit the functions to be approximated to BFs, one digit is enough (k = 1), which gives  $\Psi(0.i_1) = 0.i_1$ , *i.e.*, the identity function  $\Psi(x) = x$ . Such a solution builds simple analog neurons. They have fan-in  $\Delta \leq 2n+1$ , for which the known weight bounds (holding for any *fan-in*  $\Delta \ge 4$ ) are [52,56,61,66]:

$$2^{(\Delta-1)/2} < weight < (\Delta+1)^{(\Delta+1)/2}/2^{\Delta}$$
 (6)

Thus, a *precision* of between  $\Delta$ , and  $\Delta \log \Delta$  bits per weight would be expected. Unfortunately, the constructive solution for Kolmogorov's superpositions requires a double exponential precision for  $\psi$  (eq. 4), and for the weights:

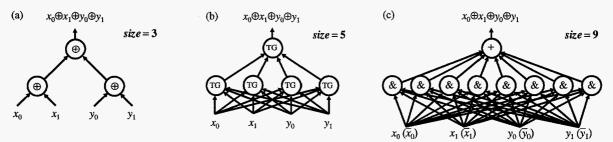


Figure 1. The PARITY problem: (a) solution using XORs gates; (b) solution using TGs; (c) solution using AND-OR gates.

$$\alpha_{p} = \sum_{r=1}^{\infty} \gamma^{-(p-1)\frac{r}{n-1}}.$$
(7)

For BFs the precision is reduced to  $(2n+2)^{-n}$ , or  $2n\log n$ bits per weight. Analog implementations are limited to just several bits [42], this being one of the reasons for investigations on precision [18,27,73,74], and on algorithms relying on limited integer weights [7,19,39]. As an example, let us consider the PARITY function of four bits. It is known that PARITY can be implemented with three 2-input XOR gates (Fig. 1.a), or with five 4-input TGs (Fig. 1.b). It is also known that a 4-input BF requires  $2\sqrt{16} + 3 = 11$  TGs [62]. A classical Boolean solution requires eight AND gates (Fig. 1.c). A brute force solution would approximate a 4input BF (Fig. 2.a and 2.b). Another solution having COM-PARISON as the  $\psi$  function, and translated inputs is presented in Fig. 2.c; it . The 2n + 1 = 5 hidden functions ( $\Phi_a$ ) are simple AND functions, while the addition is an OR function.

Due to the limited precision, an optimal solution for implementing BFs should decompose the given BF in simpler BFs (which can be implemented based on Kolmogorov's superpositions). The partial results from this first layer can be combined using (again) Kolmogorov's superpositions. Such an analog implementation will requires more than three layers. A systematic solution which would utilise silicon to the best advantage would be to rewrite a given computation (i.e., set of BFs) in a base larger than 2, and use Kolmogorov's superpositions for implementing the digitwise computations in this larger base.

#### **IV. CONCLUSIONS**

Arbitrary BFs can be implemented using: (i) Boolean circuits, but require exponential size; (ii) TGs, but (again) in exponential size; (iii) analog building blocks in linear size (having linear fan-in and polynomial precision).

The main conclusion is that size-optimal hardware im-

plementations of BFs can be obtained using analog (or mixed analog/digital) circuitry. The high precision required by the solution based on Kolmogorov's superpositions can be tackled by decomposing a BFs into simpler BFs. This is mathematically equivalent to computing in a larger base. Due to the reduced number of inputs, Kolmogorov's superpositions can be used to design the analog implementations of the digit-wise computations in such larger bases.

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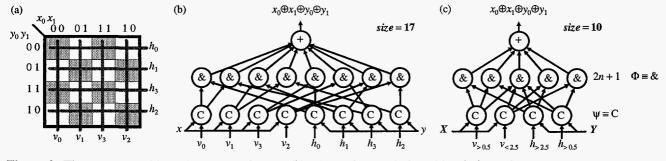


Figure 2. The PARITY problem: (a) Karnaugh map; (b) a brute force solution; (c) solution using COMPARATORS.

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